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## Total [1, k]-Sets in the Lexicographic Product of Graphs

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**Abstract.** A subset  $S \subseteq V$  in a graph G = (V, E) is called a [1, k]-set, if for every vertex  $v \in V \setminus S$ ,  $1 \leq |N_G(v) \cap S| \leq k$ . The [1, k]-domination number of G, denoted by  $\gamma_{[1,k]}(G)$  is the size of the smallest [1, k]-sets of G. A set  $S' \subseteq V(G)$  is called a total [1, k]-set, if for every vertex  $v \in V$ ,  $1 \leq |N_G(v) \cap S| \leq k$ . If a graph G has at least one total [1, k]-set then the cardinality of the smallest such set is denoted by  $\gamma_{t[1,k]}(G)$ . In this paper, we investigate the existence of [1, k]-sets in lexicographic products  $G \circ H$ . Furthermore, we completely characterize graphs whose lexicographic product has at least one total [1, k]-set. Finally, we show that finding smallest total [1, k]-set is an NP-complete problem.

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### 1. Introduction and Terminology

The concept of dominating set and domination number is a well studied topic in graph theory and has many extensions and applications [8, 9]. Many variants of domination numbers have been proposed and surveyed in the literature such as total domination number [10], efficient and open efficient domination numbers

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[1], k-tuple domination number [2] and others like [8]. Most of these problems are shown to be NP-hard. Recently, Chellali et al. introduced the notion of a [j, k]-dominating set for a graph and studied some problems in this respect [4]. They have also pointed out a number of open problems on [1, 2]-dominating sets in [4]. Some of those problems are solved by X. Yang et al. [12] and AK. Goharshady et al. [5].

In [3], Chellali et al. investigated independent [1, k]-sets for graphs and gave a constructive characterization of the trees having an independent [1, k]-set. Also they proved that the corona of two graphs G and H has an independent [1, k]-set if and only if each component of G is an isolated vertex or  $i(H) \leq k$ , where  $i(H) = i_{[1,k]}(G)$  is the minimum cardinality of an independent [1, k]-set of G.

All graphs in this paper are assumed to be a simple ones, i.e., finite, undirected, loopless graphs without multiple edges. For notation and terminology that are not defined here, we refer the reader to [11]. For given simple graph G with vertex set V(G) and edge set E(G), the degree of vertex  $v \in V(G)$  is denoted by  $d_G(v)$ , or simply d(v). We denote the minimum and maximum degrees of vertices in G by  $\delta(G)$  and  $\Delta(G)$ , respectively. The open neighborhood  $N_G(v)$ of a vertex  $v \in V(G)$  equals  $\{u : \{u, v\} \in E(G)\}$  and its closed neighborhood  $N_G[v]$  is defined  $N_G(v) \cup \{v\}$ . The open (resp. closed) neighborhood of  $S \subseteq V$  is defined to be the union of open (resp. closed) neighborhoods of vertices in S and is denoted by N(S) (resp. N[S]). A set  $D \subseteq V$  is called a dominating set of G if for every  $v \in V \setminus D$ , there exists some vertex  $u \in D$  such that  $v \in N(u)$ . The domination number of G is the minimum number among cardinalities of all dominating sets of G and is denoted by  $\gamma(G)$ . A set  $D \subseteq V$  is called a total dominating set of G if for every  $v \in V$ , there exists some vertex  $u \in D$  such that  $v \in N(u)$ . The total domination number is the minimum number among cardinalities of all total dominating sets of G and is denoted by  $\gamma_t(G)$ . For two given integers j and k such that  $j \leq k$ , a subset  $D \subseteq V$  is called a [j,k]-set (resp. total [j,k]-set) if for every vertex  $v \in V \setminus D$  (resp.  $v \in V$ ),  $j \leq |N(v) \cap D| \leq k$ . Note that total [j, k]-sets might not exist for an arbitrary graph. The sets of all graphs like G which have at least one total [j, k]-set is denoted by  $\mathcal{D}_{[i,k]}^t$ . Other types of dominating sets, that we are used in this work are summarized in the Table 1.

**Table 1:** Some types of domination studied in this paper where  $S \subseteq V$ 

Name	$v \in V \setminus S$	$v \in S$
[1,k]-set	$ N(v) \cap S  \in [1,k]$	-
Independent $[1, k]$ -set	$ N(v) \cap S  \in [1,k]$	$ N(v) \cap S  = 0$
j-dependent $[1, k]$ -set	$ N(v) \cap S  \in [1,k]$	$ N(v) \cap S  \in [0, j]$
Total $[1, k]$ -set	$ N(v) \cap S  \in [1,k]$	$ N(v) \cap S  \in [1,k]$
j-dependent total $[1, k]$ -set	$ N(v) \cap S  \in [1,k]$	$ N(v) \cap S  \in [1, j]$
Efficient dominating	$ N(v) \cap S  = 1$	$ N(v) \cap S  = 0$

The rest of the paper is organized as follows: In Section 2, we study total [1, k]sets of lexicographic product of graphs and then, we completely characterize graphs which their lexicographic product has at least one total [1, k]-set. Then, we determine the structure of all total [1, k]-sets for these graphs. In Section 3, we prove that finding a total [1, 2]-set with minimum cardinality for a graph is NP-complete.

# 2. Total [1,2]-Sets of Lexicographic Products of Graphs

The lexicographic product of graphs G and H, denoted by  $G \circ H$  is a graph with the vertex set  $V(G \circ H) = V(G) \times V(H)$  and two vertices (g, h) and (g', h') are adjacent in  $G \circ H$  if and only if either  $\{g, g'\} \in E(G)$  or g = g' and  $\{h, h'\} \in E(H)$ .

Note that if G is not connected, then  $G \circ H$  is not connected, too. So in this section, we always assume that G is a connected graph.

In this section, we investigate properties of graphs G and H such that  $G \circ H$  has a total [1, 2]-set. Then we extend these results to total [1, k]-set. Note that, it is possible that  $G \in \mathcal{D}_{[1,2]}^t$ , whereas  $G \circ H \notin \mathcal{D}_{[1,2]}^t$ , or vice versa.

**Definition 2.1.** Let H and G be two graphs. Let  $g_0 \in V(H)$  and  $h_0 \in V(H)$ . The sets  $G^{h_0} = \{(g, h_0) \in V(G \circ H) : g \in V(G)\}$  and  $H^{g_0} = \{(g_0, h) \in V(G \circ H) : h \in V(H)\}$  are called a G-Layer and a H-Layer respectively.

**Lemma 2.2.** Let v and v' be two adjacent vertices of G and  $u, u' \in V(H)$ . Then

$$N_{G \circ H}((v, u)) \cup N_{G \circ H}((v', u)) = N_{G \circ H}((v, u')) \cup N_{G \circ H}((v', u'))$$
  
=  $N_{G \circ H}((v, u)) \cup N_{G \circ H}((v', u')).$ 

**Proof.** We know that

$$N_{G \circ H}((v, u)) = \bigcup_{v_i \in N_G(v)} V(H^{v_i}) \cup \{(v, u_j) : u_j \in N_H(u)\},\$$

 $\mathbf{SO}$ 

$$N_{G\circ H}((v,u)) \cup N_{G\circ H}((v',u')) = D_1 \cup D_2 \tag{1}$$

where  $D_1 = (\bigcup_{v_i \in N_G(v)} V(H^{v_i})) \cup \{(v, u_j) : u_j \in N_H(u)\}$  and  $D_2 = (\bigcup_{v_i \in N_G(v')} V(H^{v_i})) \cup \{(v', u_j) : u_j \in N_H(u')\}.$ It is easy to see that

$$\{(v, u_j) : u_j \in N_H(u)\} \subseteq V(H^v).$$

$$(2)$$

and

$$\{(v', u_j) : u_j \in N_H(u')\} \subseteq V(H^{v'}).$$

$$(3)$$

By hypotheses  $\{v, v'\} \in E(G)$ , we have

$$V(H^{v}) \subseteq N_{G \circ H}((v', u')),$$
  

$$V(H^{v'}) \subseteq N_{G \circ H}((v, u)).$$
(4)

So by Relations 1, 2, 3 and 4, it is implied that

$$N_{G \circ H}((v, u)) \cup N_{G \circ H}((v', u')) = \bigcup_{v_i \in N_G(\{v, v'\})} V(H^{v_i}).$$

The equation above shows that the union of neighbors of the vertices (v, u) and (v', u') is independent from u and u'. Therefore, we have

$$N_{G \circ H}((v, u)) \cup N_{G \circ H}((v', u)) = N_{G \circ H}((v, u')) \cup N_{G \circ H}((v', u'))$$
  
=  $N_{G \circ H}((v, u)) \cup N_{G \circ H}((v', u')).$ 

**Lemma 2.3.** Let D be a total [1, 2]-set for  $G \circ H \in \mathcal{D}_{[1,2]}^t$  which contains more than two vertices of a  $H_-Layer H^v$ . Then  $G = K_1$  and  $H \in \mathcal{D}_{[1,2]}^t$ .

**Proof.** Let *D* be a total [1, 2]-set of  $G \circ H$  that contains vertices (x, v), (y, v) and (z, v) where  $v \in V(G)$  and  $x, y, z \in V(H)$ . If there exists a vertex  $v' \in V(G)$  such that  $\{v, v'\} \in E(G)$ , then all vertices of  $H^{v'}$  are dominated by three vertices (x, v), (y, v) and (z, v). This is a contradiction. So there is not any vertex adjacent to v. Since *G* is a connected graph,  $G = K_1 = (\{v\}, \emptyset)$  and  $S = \{u : (v, u) \in D\}$  is a total [1, 2]-set for *H* and hence  $H \in \mathcal{D}^{t}_{[1, 2]}$ .  $\Box$ 

**Corollary 2.4.** Let G be a nontrivial connected graph and  $G \circ H \in \mathcal{D}^t_{[1,2]}$ . Then, every total [1,2]-set of  $G \circ H$  has at most two vertices of each  $H_-$  Layer. For a total [1,2]-set D, we define  $A_1^D$  as  $\{(v,u) : |V(H^v) \cap D| = 1\}$  and  $A_2^D$ as  $\{(v,u) : |V(H^v) \cap D| = 2\}$ . The set D satisfies in one of the following conditions:

- 1)  $A_1^D = \emptyset$  and  $A_2^D \neq \emptyset$ ,
- 2)  $A_1^D \neq \emptyset$  and  $A_2^D \neq \emptyset$ ,
- 3)  $A_2^D = \emptyset$  and  $A_1^D \neq \emptyset$ .

**Lemma 2.5.** Let D be a total [1,2]-set of  $G \circ H \in \mathcal{D}_{[1,2]}^t$  such that  $A_2^D = \emptyset$ . Then,  $S = \{u : (u,v) \in D\}$  is a total [1,2]-set for G. In addition, if there is a vertex  $u \in S$  such that  $|N(u) \cap S| = 2$ ; then H contains an isolated vertex.

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**Proof.** The proof is by contradiction. Assume D is a total [1,2]-set of  $G \circ H$  with  $A_2^D = \emptyset$  and  $S = \{u : (u, v) \in D\}$  is not a total set of G. Then, we have three cases to consider.

- 1. There exists a vertex like  $u \in S$  such that  $|N(u) \cap S| = 0$ . It means that there is no vertex  $u' \in N_G(u)$  such that  $u' \in S$ . The set D is a total [1,2]-set and  $u \in S$ , so there exists a vertex  $v \in V(H)$  such that  $(u,v) \in D$ . Similarly there exists a vertex  $v' \in V(H)$  such that  $(u,v') \in D$ . This is a contradiction against  $A_2^D = \emptyset$ .
- 2. There exists a vertex like  $w \in V(G) \setminus S$  such that  $|N_G(w) \cap S| = 0$ . Then, there is no vertex like  $v \in V(H)$  such that  $(u, v) \in D$ . Moreover, there is no vertex  $w' \in N_G(w)$  such that  $w' \in S$ . Therefore vertices of  $H^w$  can not be dominated by any vertex in D, which is a contradiction.
- 3. There exists a vertex like  $w \in V(G) \setminus S$  such that  $|N(w) \cap S| > 2$ . Then, there are at least three distinct vertices  $w', w'', w''' \in N_G(w) \cap S$ . By the definition of S, there are vertices  $v', v'', v''' \in V(H)$  such that  $(w', v'), (w'', v''), (w''', v''') \in D$ . These vertices dominate all vertices of  $H^w$ , which is a contradiction.  $\Box$

**Lemma 2.6.** Let  $G \circ H \in \mathcal{D}_{[1,2]}^t$  and H does not contain any isolated vertex. Then, there exists either a 1-dependent total [1,2]-set for G or for each total [1,2]-set D of G,  $A_1^D = \{(v,u) : |V(H^v) \cap D| = 1\} \neq \emptyset$  and  $A_2^D = \{(v,u) : |V(H^v) \cap D| = 2\} \neq \emptyset$ .

**Proof.** Let *D* be a total [1, 2]-set of  $G \circ H$  which contains at most one vertex from each  $H_{-}$ Layer. Since *H* does not contain any isolated vertex then by Lemma 2.5 there is a 1-dependent total [1, 2]-set like *S* for *G* such that S = $\{v : (v, u) \in D\}$  and  $A_2^D = \emptyset$ .  $\Box$ 

For a given graph  $G \circ H \in \mathcal{D}_{[1,2]}^t$  and a total [1,2]-set D of  $G \circ H$  where  $A_2^D \neq \emptyset$ , we define the set  $B^D$  as  $B^D = \{\{u', u''\} : (v, u'), (v, u'') \in A_2^D\}.$ 

**Lemma 2.7.** Let  $G \circ H \in \mathcal{D}^t_{[1,2]}$  where H does not contain any isolated vertex and for any total [1,2]-set D of  $G \circ H$ ,  $A_1^D \neq \emptyset$  and  $A_2^D \neq \emptyset$ . Then, the following conditions hold:

- 1) Every element of  $B^D$  is a total [1,2]-set for H.
- 2) The set  $S' = \{v : (v, u) \in D\}$  is a 1-dependent [1, 2]-set for G.
- 3) If there is a vertex  $v \in S'$  such that  $|N(v) \cap S'| = 0$  then  $dist_G(v, v') \ge 3$ for every  $v' \in S' \setminus \{v\}$ .

**Proof.** Let *D* be a total [1,2]-set of  $G \circ H \in \mathcal{D}^t_{[1,2]}$ ; there are three cases to consider.

- 1) Suppose that  $S = \{u^*, u^{\bullet}\} \in B$  is not a total [1, 2]-set for H. Then two cases occur and in each case, we can establish a contradiction with D is a total [1, 2]-set.
  - Let  $\{u^*, u^{\bullet}\} \notin E(H)$  and there is a  $(v', u') \in D$  such that  $\{(v, u^*), (v', u')\} \in E(G \circ H)$ . Since H dose not contain any isolated vertex, so any vertex  $u'' \in N_H(u')$  is dominated by  $(v', u'), (v, u^*)$ and  $(v, u^{\bullet})$ .
  - Let  $\{u^*, u^{\bullet}\}$  does not dominate all vertices of V(H). So, there is a vertex  $(v', u') \in D$  such that  $\{v, v'\} \in E(G)$  and (v', u') dominates all vertices of  $H^v$ . Then any vertex  $u'' \in N_H(u')$  is dominated by  $(v', u'), (v, u^*)$  and  $(v, u^{\bullet})$ .
- 2) Suppose that  $S' = \{v : (v, u) \in D\}$  is not a 1-dependent [1,2]-set for G. Then, three cases occur and in each case, we have a contradiction with D being a total [1,2]-set.
  - There is a vertex  $v \in S'$  that is dominated by at least two vertices  $v', v'' \in S'$ . So there are vertices  $u, u', u'' \in V(H)$  such that  $(v, u), (v', u'), (v'', u'') \in D$ . Since H dose not contain any isolated vertex, there is a vertex  $u''' \in V(H)$  such that  $\{u, u'''\} \in E(H)$ . Then, (v, u'') is dominated by (v, u), (v', u') and (v'', u'').
  - There is a vertex  $v \in V(G) \setminus S'$  such that  $|N_G(v) \cap S'| = 0$ . So no vertex of  $H^v$  is dominated by D.
  - There is a vertex  $v \in V(G) \setminus S'$  such that  $|N_G(x) \cap S'| > 2$ . Then there are at least three vertices distinct  $v', v'', v''' \in S'$  to dominate v. By definition of S', there are vertices  $u', u'', u''' \in V(H)$  such that  $(v', u'), (v'', u''), (v''', u''') \in D$ . These vertices dominate all vertices of  $H^v$ .
- 3) Let  $v \in S'$  such that  $|N(v) \cap S'| = 0$  and there is a vertex  $v' \in S'$  such that  $dist_G(v, v') = 2$ .

By  $|N(v) \cap S'| = 0$ , there exist vertices  $u', u'' \in V(H)$  such that  $(v, u'), (v, u'') \in D$  and  $\{u', u''\} \in E(H)$ . Suppose there is a vertex  $v' \in S'$  such that  $dist_G(v, v') = 2$ . So, there is a vertex  $v'' \in V(G)$  such that  $\{v, v''\}, \{v', v''\} \in E(G)$ . The vertices (v, u'), (v, u'') and (v', u') dominate all vertices of  $H^{v''}$ . It is contradictory with D being a total [1,2]-set. So we have  $dist_G(v, v') \ge 3$ .  $\Box$ 

**Lemma 2.8.** Let D be a total [1, 2]-set of  $G \circ H \in \mathcal{D}_{[1,2]}^t$  such that  $A_1^D = \emptyset$ . Then  $S' = \{v : (v, u) \in D\}$  is an efficient dominating set of G.

**Proof.** Since *D* be a total [1,2]-set of  $G \circ H$ , then there is a vertex  $v \in S'$  such that the set *D* contains (v, u'), (v, u'') for some vertex  $u', u'' \in V(H)$ . By Lemma 2.7,  $\{u', u''\}$  is a total [1,2]-set for *H*. So for any vertex  $v' \in N_G(v)$ , none of vertices in  $H^{v'}$  cannot be contained in *D*. Thus  $dist_G(v, v') \ge 3$  and *S* is an efficient dominating set of *G*.  $\Box$ 

In the sequel  $\mathcal{SD}_{[i,j]}^k(G)$  is used to denote the set of all k-dependent [i,j]-set S of G such that S satisfies in the following condition

$$(\forall v \in S | N(v) \cap S| = 0) \to (\forall v' \in S \setminus \{v\} d(v, v') \ge 3).$$

**Corollary 2.9.** Let G be a connected nontrivial graph and D be a total [1, 2]-set of  $G \circ H \in \mathcal{D}_{[1,2]}^t$ , one of the following cases holds:

- If  $A_1^D = \{(u, v) : |V(H^v) \cap D| = 1\} = \emptyset$ , then there is a total [1,2]-set  $S = \{u^*, u^\bullet\}$  in H and an efficient dominating set S' in G such that  $D' = S' \times S$  is a total [1,2]-set for  $G \circ H$  and |D| = |D'| = 2|S'|.
- If  $A_2^D = \{(u,v) : |V(H^v) \cap D| = 2\} = \emptyset$  and H contains an isolated vertex v. Then there is a total [1,2]-set S in G where  $D' = S \times \{v\}$  and D' is a total [1,2]-set for  $G \circ H$ . Moreover, we have |D| = |D'| = |S|.
- If  $A_2^D = \{(u, v) : |V(H^v) \cap D| = 2\} = \emptyset$  and H does not contain any isolated vertex, then for every vertex  $v \in V(H)$  there is a 1-dependent total [1, 2]-set S in G such that  $D' = S \times \{v\}$  and D' is a total [1, 2]-set for  $G \circ H$ . Clearly, |D| = |D'| = |S|.
- If  $A_1^D \neq \emptyset$  and  $A_2^D \neq \emptyset$ , then there is a total [1,2]-set  $S = \{u^*, u^*\}$  in Hand a 1-dependent total [1,2]-set S' in G such that for any vertex  $v \in S$ and  $u \in X$  where  $X = \{x : |N_G(x) \cap S'| = 0\}$ ,  $dist(v, u) \ge 3$ . Moreover  $D' = ((X \times S) \cup (S' \setminus X) \times \{u^*\})$  is a total [1,2]-set of size |D| in  $G \circ H$ and |D| = |D'| = |S'| + |X|.

**Proof.** This corollary is a direct result of Lemma 2.2, 2.5, 2.7 and 2.8.  $\Box$ 

**Theorem 2.10.** Let G and H be two graphs. Then,  $G \circ H \in \mathcal{D}_{[1,2]}^t$  if and only if one of the following conditions holds:

- 1.  $G = K_1$  and  $H \in \mathcal{D}^t_{[1,2]}$ ;
- 2. G has a total [1,2]-set S such that if S has a vertex v where  $|N(v) \cap S| = 2$ then H has an isolated vertex;
- 3. G is an efficient domination graph and  $\gamma_{t[1,2]}(H) = 2$ ;

4.  $SD^{1}_{[1,2]}(G) \neq \emptyset$  and  $\gamma_{t[1,2]}(H) = 2$ .

**Proof.** Suppose that D be a total [1, 2]-set of  $G \circ H \in \mathcal{D}^t_{[1,2]}$ . If D contains more than two vertices of a  $H_-$ Layer, then by Lemma 2.3,  $G = K_1$  and  $H \in \mathcal{D}^t_{[1,2]}$ . If D contains at most two vertices of each  $H_-$ Layer, then there is a total [1, 2]-set D' for  $G \circ H$  such that |D'| = |D| and vertices of D' have been choosen from two  $G_-$ Layers as  $G^{u^*}$  and  $G^{u^\bullet}$ . Without lose of generality we consider that  $S = \{v : (v, u) \in D'\}$  and  $S' = \{u^*, u^\bullet\}$ . Then, the set D' satisfies one of the following conditions:

- a) By Lemma 2.5,  $D = \{(v, u^*) : v \in S\}$ , so S is a total [1, 2]-set for G and if there exists a vertex  $v \in D$  such that  $|N(v) \cap S| = 2$ , then H has an isolated vertex.
- b)  $D' = \{(v, u^*) : v \in S \text{ and } u \in S'\}$ , by Corollary 2.9, S is an efficient dominating set of G and S' is a total [1, 2]-set for H.
- c) There is a vertex  $w \in S$  such that  $(w, u^*) \in D'$  but  $(w, u^{\bullet}) \notin D'$ . By Lemma 2.7, we have  $S \in SD^1_{[1,2]}(G)$  and S' is a total [1,2]-set for H.

Now, we show the other side as follows:

- 1. If  $G = K_1$  and H has a total [1,2]-set S', then it is easy to see that  $G \circ H = H$  and S' is a total [1,2]-set of  $G \circ H$ .
- 2. Assume that S is a total [1,2]-set of G and  $u^* \in V(H)$ . We define D as  $S \times \{u^*\}$ . Since every vertex of  $G^{u^*}$  is dominated by at least one of vertices of D, then every vertex of other  $G_{-}$ Layers is dominated by D. So, for any vertex  $(v', u') \in G \circ H$ , we have  $|N((v', u')) \cap D| \ge 1$ . Now, it is sufficient to show that  $|N((v', u')) \cap D| \le 2$ . To this end, we consider two cases:
  - a) For every vertex  $v \in S$ ,  $|N(v) \cap S| = 1$ : So, it is clear that for any vertex  $(v', u^*)$  of  $G^{u^*}$ ,  $|N((v', u^*)) \cap D| \leq 2$ . If  $u' \neq u^*$ , we need to show that  $|N((v', u')) \cap D| \leq 2$ . Then following cases can happen:
    - a1)  $(v', u^*) \in D$  and  $\{u', u^*\} \in E(H)$ ; for every  $v'' \in S$  adjacent to v', (v', u') is dominated by  $(v', u^*)$  and  $(v'', u^*)$ . Since  $(v', u^*) \in D$  and  $v' \in S$ , so  $|N(v') \cap S| = 2$  and  $|N((v', u')) \cap D| = |N(v') \cap S| + 1 = 2$ .
    - a2)  $(v', u^{\star}) \in D$  and  $\{u', u^{\star}\} \notin E(H)$ ; if  $v'' \in S$  and  $\{v', v''\} \in E(G)$  then (v', u') is dominated by  $(v'', u^{\star})$ . So  $|N((v', u')) \cap D| = |N(v') \cap S| = 1$ .

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- a3)  $(v', u^{\star}) \notin D$ ; for every  $v'' \in S$  and  $\{v', v''\} \in E(G)$ , (v', u')is dominated by  $(v'', u^{\star})$ . Since  $(v', u^{\star}) \notin D$ ,  $v' \notin S$ . We have $|N((v', u')) \cap D| = |N(v') \cap S| \leq 2$ .
- b) There is a vertex  $v \in S$  such that  $|N(v) \cap S| = 2$  and  $u^*$  is an isolated vertex in H. For every vertex  $v'' \in S$  and  $\{v', v''\} \in E(G), (v', u')$  is dominated by  $(v'', u^*)$ . So it is the case that  $|N((v', u')) \cap D| = |N(v') \cap S| \leq 2$ .
- 3. Let S be an efficient dominating set of G,  $S' = \{u^*, u^\bullet\}$  is a total [1,2]set for H and  $D = \{(v, u) : v \in S \text{ and } u \in S'\}$ . It is easy to see that D is a total dominating set of  $G \circ H$ .

If  $v' \in S$ , then every  $(v', u') \in V(H^{v'})$  are dominated by either  $(v', u^*)$ or  $(v', u^{\bullet})$ . Since S is an efficient dominating set of G, then  $N_G(v') \cap S = \emptyset$  and (v', u') is not dominated by any other vertices. If  $v' \notin S$ , then there is exactly one vertex  $v'' \in S$  such that  $\{v', v''\} \in E(G)$  and every  $(v', u') \in V(H^{v'})$  are dominated by either  $(v'', u^*)$  and  $(v'', u^{\bullet})$ . So, D is a total [1, 2]-set for  $G \circ H$ .

4. Suppose that  $S \in \mathcal{SD}^1_{[1,2]}, S' = \{u^\star, u^\bullet\}$  is a total [1,2]-set for H and

$$D = \{(v, u^*), (v, u^\bullet) : v \in S \text{ and } |N(v) \cap S| = 0\}$$
$$\cup \{(v, u^*) : v \in S \text{ and } |N(v) \cap S| = 1\}.$$

By definition of D, It is easy to see that for any vertex  $(v, u) \in D$ , there is a vertex  $(v', u') \in D$  such that  $\{(v, u), (v', u')\} \in E(G \circ H)$ . So, D is a total set of  $G \circ H$ . Now, we must show that D dominates all vertices of  $G \circ H$  at least one and at most two times. It is clear  $S = \{v : (v, u^*) \in D\} \in S\mathcal{D}^1_{[1,2]}$ . We consider three kind of vertices and we will show vertices of each  $H_-$ Layer are dominated by at least one and two vertices of D.

- a)  $v \in S$  and  $|N(v) \cap S| = 0$ : Since  $S' = \{u^*, u^\bullet\}$  is a total [1,2]-set for  $G \circ H$ ,  $(v, u^*) \in D$  and  $(v, u^\bullet) \in D$ . Then, all of the vertices of  $H^v$  are dominated by  $(v, u^*)$  and  $(v, u^\bullet)$ . Since  $|N(v) \cap S| = 0$ . So, any other vertex cannot dominate vertices of  $H^v$ . Therefore  $1 \leq |N(v, u) \cap D| \leq 2$ .
- b)  $v \in S$  and  $|N(v) \cap S| = 1$ : So, there is a vertex  $v' \in S$  such that  $\{v, v'\} \in E(G)$ ,  $(v', u^*)$  dominates all of the vertices of  $H^v$  and these vertices can also be dominated by  $(v, u^*)$ . Since S is a 1-dependent [1, 2]-set for G, then there is not any other vertex in neighborhood of v in S, so  $1 \leq |N(v, u) \cap D| \leq 2$ .

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c)  $v \notin S$ : Since S is a 1-dependent [1,2]-set for G, it is easy to see that there is a vertex  $v' \in S$  such that  $\{v, v'\} \in E(G)$ . So, all of the vertices of  $H^v$  are dominated by  $(v', u^*)$ . If  $|N(v') \cap S| = 0$ , then  $(v', u^{\bullet})$  dominates vertices of  $H^v$  and any other vertices can not dominate them. If there exist a  $v'' \in S$  such that  $\{v, v''\} \in E(G)$ and it is contradict to  $dist_G(v', v'') \geq 3$ . If  $|N(v') \cap S| = 0$ , there maybe exists a vertex  $(v'', u^*) \in D$  such that  $|N(v') \cap S| \neq 0$  and there is no vertex in  $H^{v''}$  and other  $H_{-}$ Layers dominate vertices of  $H^v$ .  $\Box$ 

In the sequel, we express necessary and sufficient conditions for the given graphs G and H such that  $G \circ H$  has a total [1, k]-set. The Lemma 2.3, 2.5, 2.7 and Corollary 2.9 are generalized to total [1, k]-set. Since proofs in this section can be similarly obtained from the case on total [1, 2]-sets, we omit them.

**Theorem 2.11.** Let D be a total [1, k]-set for  $G \circ H$ .

- a) If D contains more than k vertices of a  $H_{-}Layer$ , then  $G = K_1$  and  $H \in \mathcal{D}_{[1,k]}^t$ .
- b) If D contains at most one vertex of every  $H_{-}Layers$ , then  $S = \{v \in V(G) : (v, u) \in D\}$  is a (k 1)-dependent total [1, k]-set of G. Moreover if there is a vertex  $v \in S$  such that  $|N(v) \cap S| = k$ , then H contains an isolated vertex.
- c) If H does not contain any isolated vertex and  $S = \{v \in V(G) : (v, u) \in D\}$  is not a total set of G, then D contains at most k vertices of each  $H^v$  and satisfies the following conditions:
  - c1) The set  $S' = \{u \in V(H) : (v, u) \in D\}$  is a total [1, k]-set of H with cardinality to at most k and there is a vertex  $x \in S$  such that  $1 < |D \cap V(H^x)| \leq |S'|;$
  - c2) S is a (k-1)-dependent [1,k]-set for G;
  - c3) If there exist a vertex  $v \in S$  such that  $|N(v) \cap S| = 0$ , then  $1 < |D \cap V(H^v)| \leq \lfloor k/2 \rfloor$  or for any vertex  $v' \in S \{v\}$ , we have  $dist_G(v, v') \geq 3$ .

**Theorem 2.12.** Let G and H be two graphs.  $G \circ H \in \mathcal{D}_{[1,k]}^t$  if and only if G and H satisfy one of the following conditions

- 1.  $G = K_1 \text{ and } H \in \mathcal{D}^t_{[1,k]};$
- 2. G has a total [1, k]-set S and if S has a vertex v such that  $|N(v) \cap S| = k$  then H has an isolated vertex;

- 3. G is an efficient domination graph and  $\gamma_{t[1,k]}(H) \leq k$ ;
- 4. G has a (k-1)-dependent [1,k]-set S and if  $S \in SD^{k-1}_{[1,k]}(G)$  then  $\gamma_{t[1,k]}(H) \leq k$  and otherwise  $\gamma_{t[1,k]}(H) \leq k/2$ .

### 3. Complexity

In this section, we will show that the decision problem for total [1, 2]-set is NP-complete. We will do this by reduction the NP-complete problem, Exact 3-Cover, to Total [1, 2]-Set.

**Exact 3-cover problem:** Input of this problem is a finite set  $X = \{x_1, x_2, ..., x_{3q}\}$ 

with |X| = 3q and a collection C of 3-element subsets of X such as  $C_i = \{x_{i_1}, x_{i_2}, x_{i_3}\}$ . our goal is to understand is there a  $C' \subseteq C$  such that every element of X appears in exactly one element of C'?

**Total** [1,2]-set problem: Input of this problem is a graph G = (V, E) and a positive integer  $k \leq |V|$ . We want to investigate is there any total [1,2]-set of cardinality at most k for G.

**Theorem 3.1.** Total [1, 2]-SET is *NP*-complete for bipartite graphs.

**Proof.** Let  $D \subseteq V$  is given, we verify D is a total [1, 2]-set. For any vertex  $v \in D$ , we check neighborhood of each vertex and compute span number of any vertex  $v \in V$ . If there is a vertex v with span number more than 2, this set isn't a total [1, 2]-set for G. It is obvious this algorithm is done in polynomial time and total [1, 2]-set is a NP problem. Now for a set X, and a collection C of 3-element subsets of X, we build a graph and transform EX-ACT 3-COVER into a total [1, 2]-set problem. Let  $X = \{x_1, x_2, ..., x_{3q}\}$  and  $C = \{C_1, C_2, ..., C_t\}$ . For each  $C_i \in C$ , we build a cycle  $C_4$  with a vertex  $u_i$  and add new vertices  $\{v_{1_1}, v_{1_2}, v_{1_3}, v_{2_1}, v_{2_3}, \cdots, v_{t_1}, v_{t_2}, v_{t_3}\}$ . We and connect all vertices  $v_{i1}, v_{i2}, v_{i3}$  to  $u_i$ . Then add some other vertices  $\{x_1, x_2, ..., x_{3q}\}$  and edges  $x_i v_{j_1}, x_i v_{j_2}$  and  $x_i v_{j_3}$ , if  $x_i \in C_j.G$  is a bipartite graph.

Let k = 2t + q. Suppose that C' is a solution for set X and collection C of EXACT 3-COVER. We build a set D of vertices of G contain every  $u_i$ ,  $1 \leq i \leq t$ , and another vertex of  $C_4$  adjacent to  $u_i$  and one of the  $v_{j_1}, v_{j_2}$  or  $v_{j_3}$  for each  $C_j \in C'$ . If C' exists, then it's cardinality is precisely q, and so |D| = 2t + q = k. We can check easily that D is a [1, 2]-total set of G.

Conversely, suppose that G has a total [1, 2]-set D with  $|D| \leq 2t + q = k$ . Then D must contain two vertices of every  $C_4$ , in the best case we select  $u_i$  and one of the vertices in that adjacency in  $C_4$ . We select 2t vertices that dominate all

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vertices of cycles and all vertices of form  $v_{i_1}, v_{i_2}$  or  $v_{i_3}$  for  $1 \leq i \leq t$ . Since each  $v_{i_j}$  dominates only three vertices of  $\{x_1, x_2, ..., x_{3q}\}$  We have to select exactly q vertices of them, i.e. we select q 3-element subsets of form  $\{v_{i_1}, v_{i_2}, v_{i_3}\}$  and one element of each of them. Each of this  $v_{i_j}$  correspond to a  $C_i$  and union of them is a exact cover for C.  $\Box$ 

**Example 3.2.** Let  $C = \{C_1, C_2, C_3, C_4\}$  where  $C_1 = \{x_1, x_2, x_4\}, C_2 = \{x_3, x_5, x_7\}, C_3 = \{x_4, x_5, x_6, x_7\}$  and  $C_4 = \{x_6, x_8, x_9\}$ , Corresponding graph was shown in Figure 1.



Figure 1. NP-completeness for bipartite graph

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